

# Overview of Intelligent RAM (IRAM)

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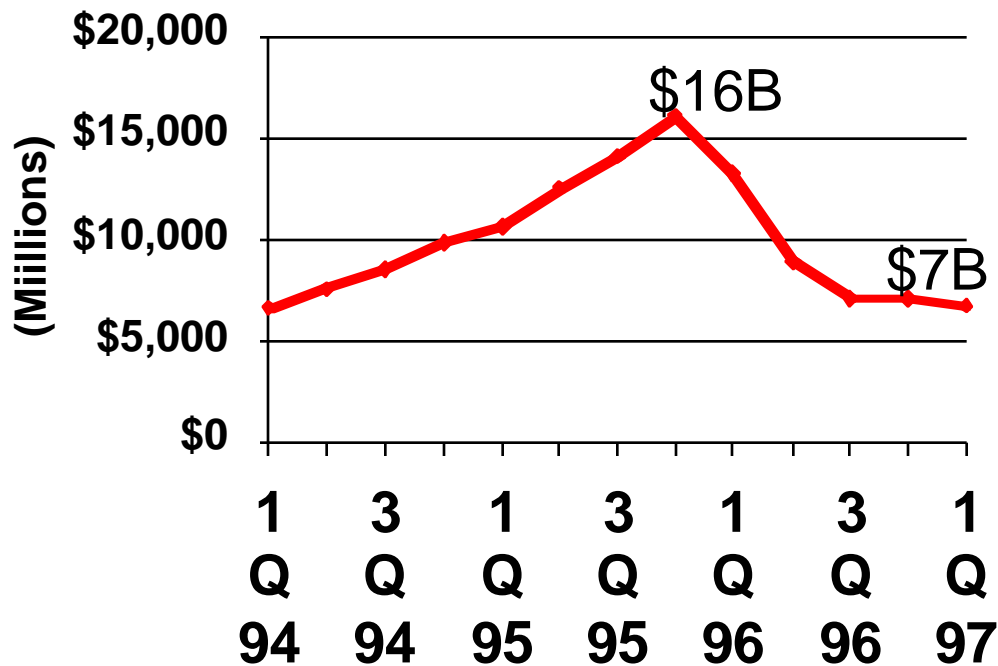
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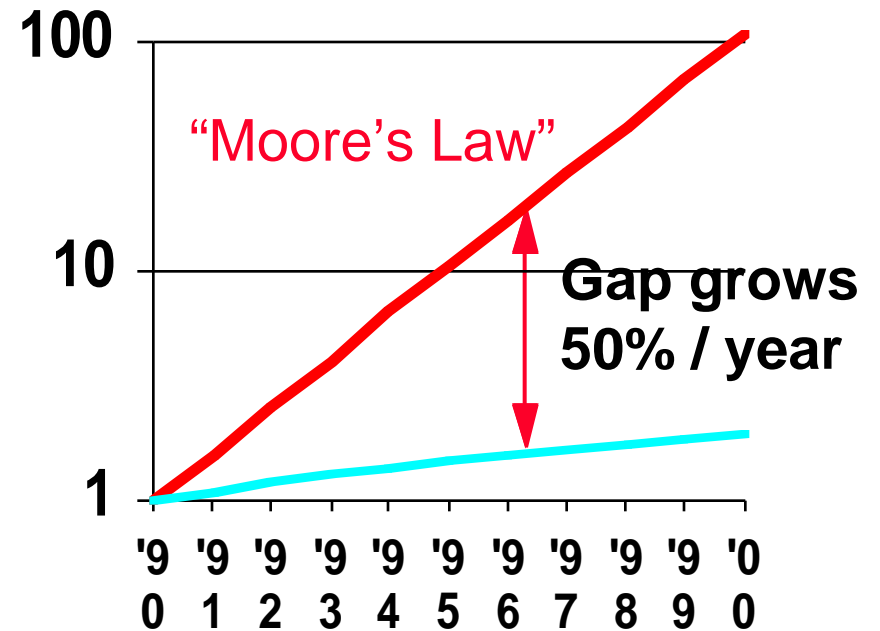
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# Motivations for IRAM

**DRAM Sales per Quarter**



**Microprocessor-DRAM Performance Gap**

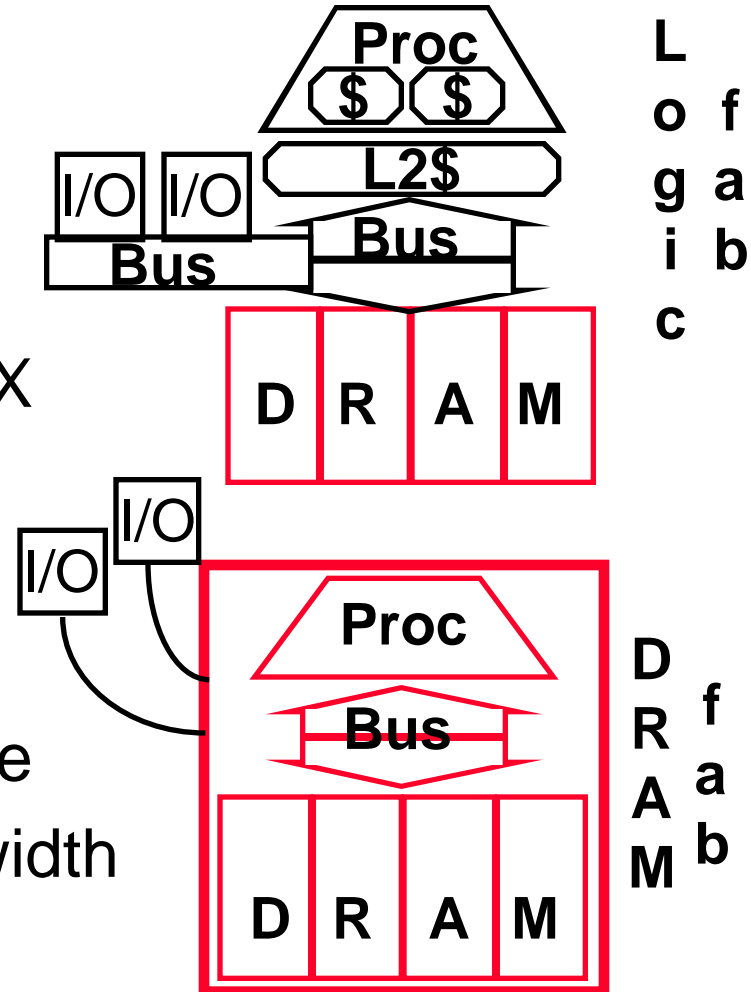


- More profit for DRAM industry?
- MPU close gap?

# IRAM Vision Statement

Microprocessor & DRAM  
on a single chip:

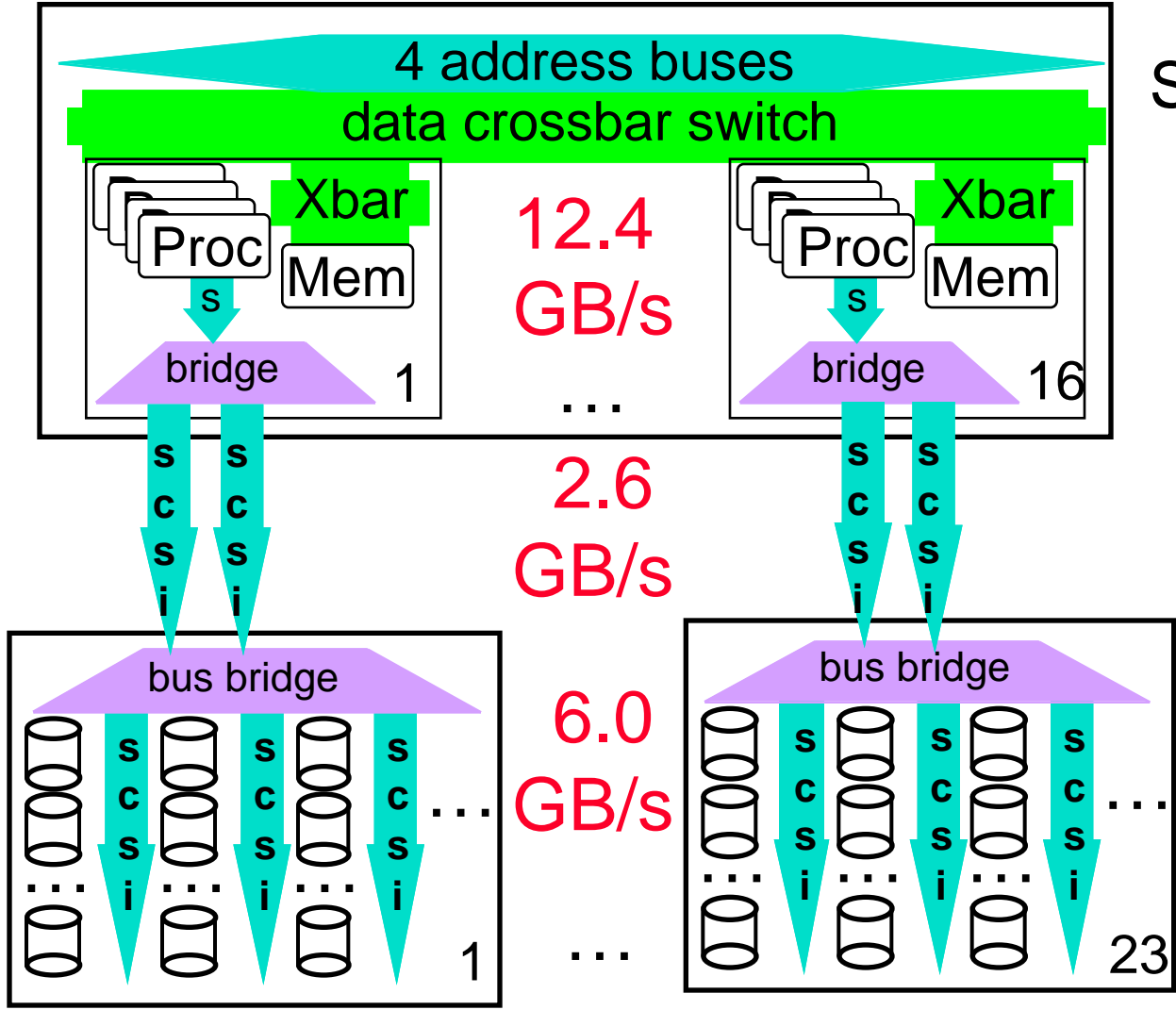
- on-chip memory latency 5-10X, bandwidth 50-100X
- improve energy efficiency 2X-4X (no off-chip bus)
- serial I/O 5-10X v. buses
- smaller board area/volume
- adjustable memory size/width



# Near-term IRAM Applications

- “Intelligent” Set-top
  - 2.6M Nintendo 64 ( $\approx$  \$150) sold in 1st year
  - 4-chip Nintendo  $\Rightarrow$  1-chip: 3D graphics, sound, fun!
- “Intelligent” Personal Digital Assistant
  - 1.0M PalmPilots ( $\approx$  \$300) sold in 1st year
  - Super PDA/Smart Phone: speech I/O + “voice” email...

# Long-term App: Decision Support?

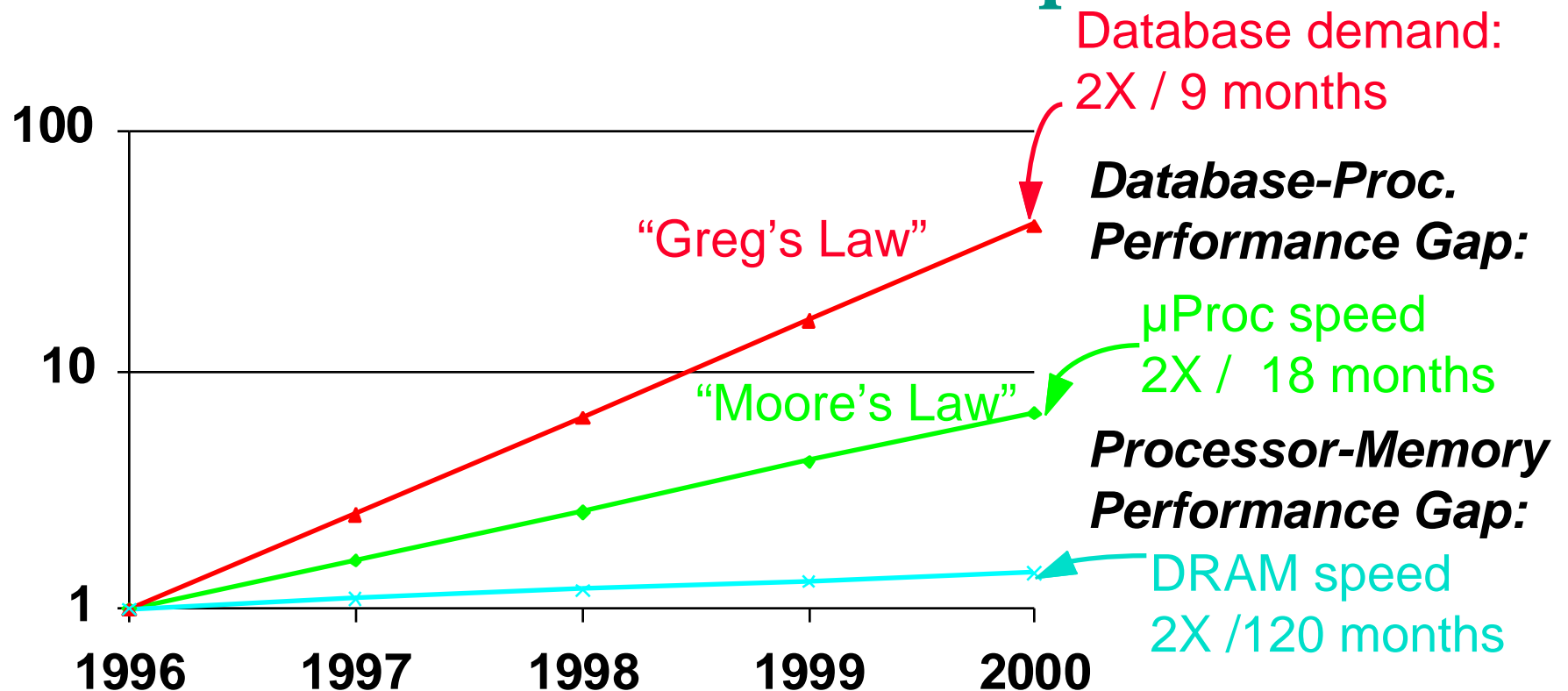


Sun 10000 (Oracle 8):

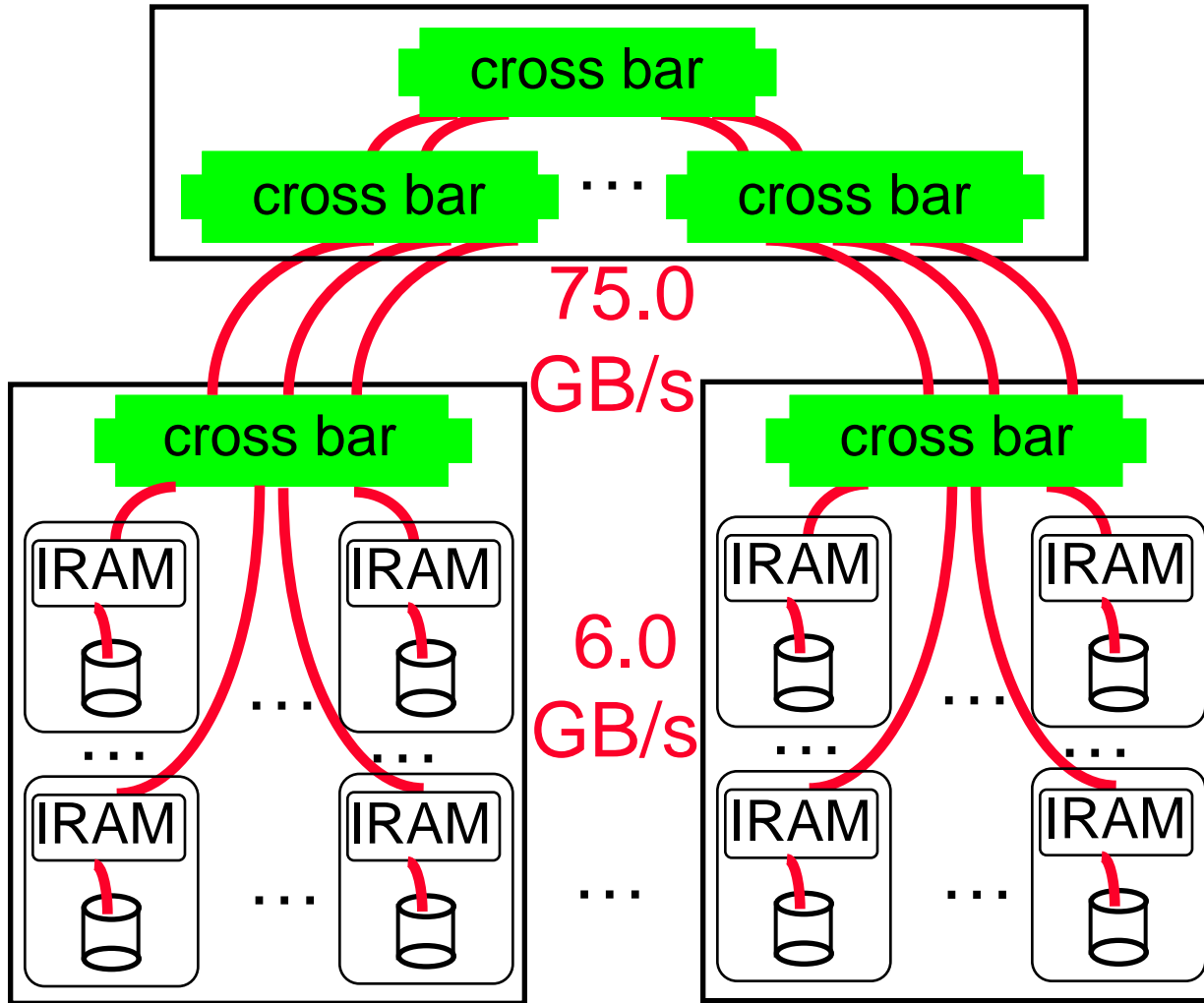
- TPC-D (1TB) leader
- SMP 64 CPUs, 64GB dram, 603 disks

Disks, encl.	\$2,348k
DRAM	\$2,328k
Boards, encl.	\$983k
CPUs	\$912k
Cables, I/O	\$139k
Misc	\$65k
HW total	<u>\$6,775k</u>

# IRAM Application Inspiration: Database Demand vs. Processor/DRAM speed



# “Intelligent Disk”: Scalable Decision Support?



1 IRAM/disk + shared nothing database

- 603 CPUs,  
14GB dram, 603 disks

Disks (market) \$840k

IRAM (@\$150) \$90k

Disk encl., racks \$150k

Switches/cables \$150k

Misc \$60k

Subtotal \$1,300k

Markup 2X? ≈ \$2,600k

≈ 1/3 price, 2X-5X perf<sub>7</sub>

# New Architecture Directions

- “...wires are not keeping pace with scaling of other features. ... In fact, for CMOS processes below 0.25 micron ... *an unacceptably small percentage of the die will be reachable during a single clock cycle.*”
- “Architectures that require long-distance, rapid interaction will not scale well ...”
  - “Will Physical Scalability Sabotage Performance Gains?” Matzke, *IEEE Computer* (9/97)



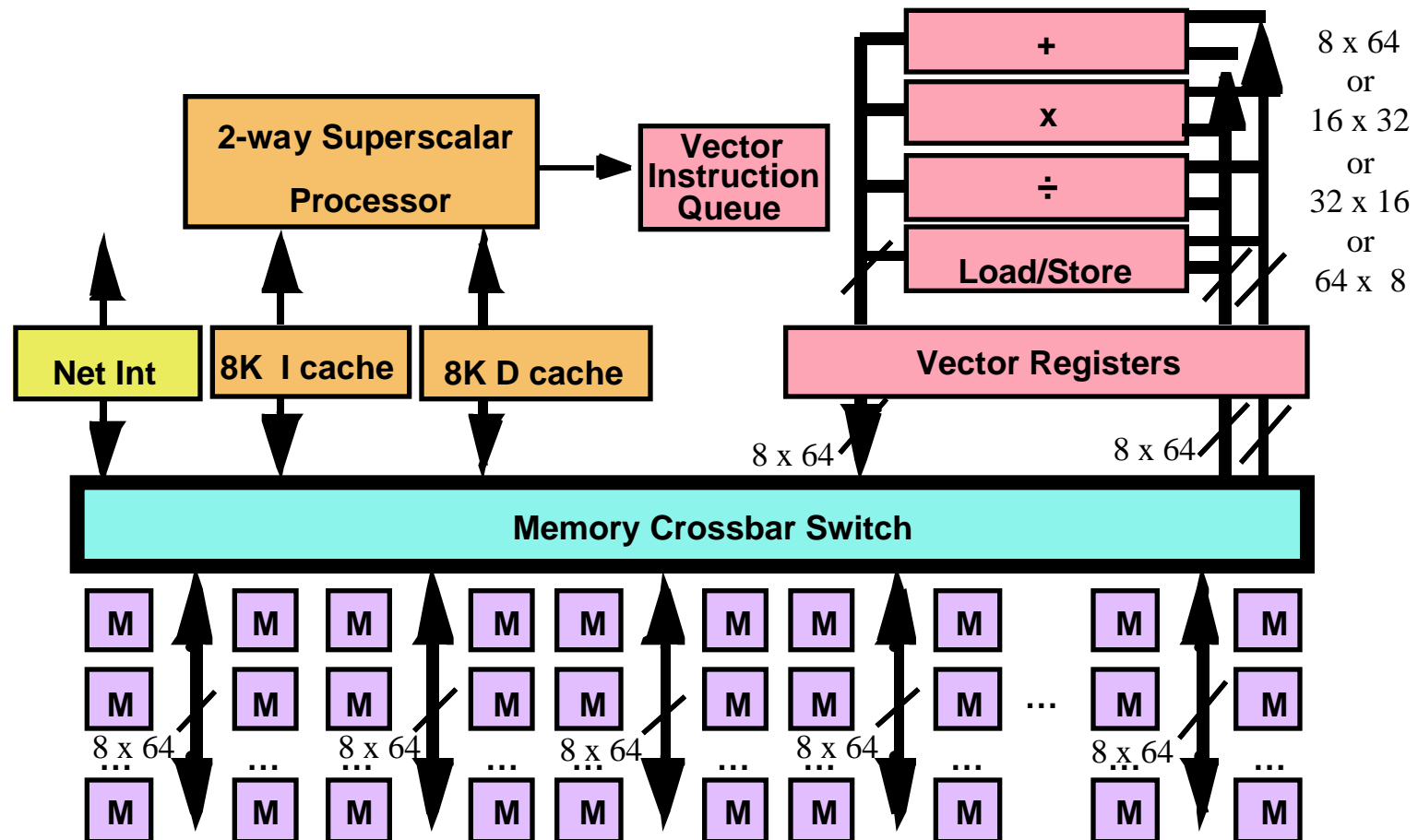
# New Architecture Directions

- “...media processing will become the dominant force in computer arch. & microprocessor design.”
- “... new media-rich applications... involve significant real-time processing of continuous media streams, and make heavy use of vectors of packed 8-, 16-, and 32-bit integer and Fl. Pt.”
- Needs include high memory BW, high network BW, continuous media data types, real-time response, fine grain parallelism
  - “How Multimedia Workloads Will Change Processor Design”, Difendorff & Dubey, *IEEE Computer* (9/97)

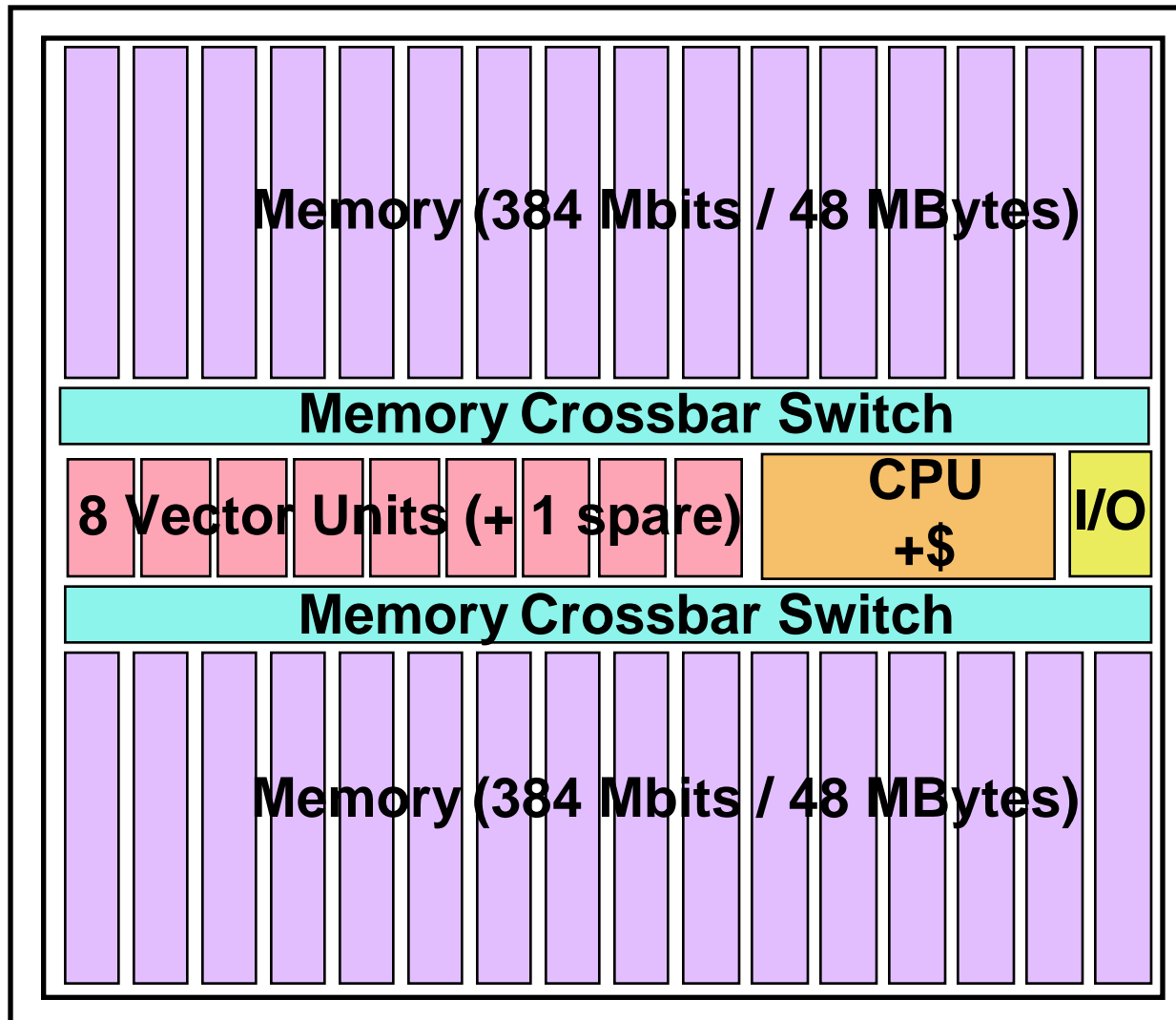
# Revive Vector Architecture!

- High cost:  
≈ \$1M / processor?
- Low latency, high BW memory system?
- Compilers?
- Performance?
- Limited to scientific applications?
- Real-time?
- Single-chip CMOS microprocessor/IRAM
- IRAM = low latency, high bandwidth memory
- For sale, mature (>20 years)
- Easy to scale speed with technology (e.g, hides latency)
- Multimedia apps vectorizable too:  
 $N*64b, 2N*32b, 4N*16b, 8N*8b$
- No caches, no speculation  
⇒ repeatable speed as vary input

# V-IRAM-2: 0.13 $\mu\text{m}$ , Fast Logic, 1GHz 16 GFLOPS(64b) / 128 GOPS(8b) / 96MB



# V-IRAM-2 Floorplan



- 0.13  $\mu\text{m}$ ,  
1 Gbit DRAM
- 1B Xtors:  
90% Memory,  
Xbar, Vector  
 $\Rightarrow$  **regular  
design**
- Spare VU &  
Memory  $\Rightarrow$   
**90% die  
repairable**
- Short signal  
distance  $\Rightarrow$   
**speed scales  
<0.2  $\mu\text{m}$**

# IRAM Conclusion

- IRAM potential in bandwidth (memory and I/O), latency, energy, capacity, board area; challenges in power/performance, testing, yield
- V-IRAM can show potential (+compilers,+testing)
- 10X-100X improvements based on technology shipping for 20 years (not JJ, photons, MEMS, ...)
- Potential upheaval in database server industry?
- Potential shift in balance of power in DRAM/microprocessor industry in 5-7 years?
  - Who ships the most memory?
  - Who ships the most microprocessors?

# Interested in Participating?

- Looking for ideas of IRAM enabled apps
- Contact us if you're interested:  
`http://iram.cs.berkeley.edu/`  
`email: patterson@cs.berkeley.edu`
- Thanks for advice/support: DARPA, Intel, LG Semiconductor, Neomagic, Samsung, SGI/Cray, Sun Microsystems

# Backup Slides

*(The following slides are used to help answer questions)*

# Characterizing IRAM

## Cost/Performance

- Cost  $\approx$  embedded processor + memory
- Small memory on-chip (25 - 100 MB)
- High vector performance (2 -16 GFLOPS)
- High multimedia performance (4 - 64 GOPS)
- Low latency main memory (15 - 30ns)
- High BW main memory (50 - 200 GB/sec)
- High BW I/O (0.5 - 2 GB/sec via N serial lines)
  - Integrated CPU/cache/memory with high memory BW ideal for fast serial I/O



# IRAM Challenges

## ■ Chip

- Good performance and reasonable power?
- Speed, area, power, yield, cost in DRAM process?
- BW/Latency oriented DRAM tradeoffs?
- Testing time of IRAM vs DRAM vs microprocessor?
- Reconfigurable logic to make IRAM more generic?

## ■ Architecture

- How to turn high memory bandwidth into performance for real applications?
- Extensible IRAM: Large program/data solution? (e.g., external DRAM, clusters, CC-NUMA, ...)

# Why IRAM now?

## Lower risk than before

- Faster Logic + DRAM available now/soon?
- DRAM manufacturers now willing to listen
  - Before not interested, so early IRAM = SRAM
- Past efforts memory limited  $\Rightarrow$  multiple chips
  - $\Rightarrow$  1st solve the unsolved (parallel processing)
    - Gigabit DRAM  $\Rightarrow$   $\approx$ 100 MB; OK for many apps?
- Systems headed to 2 chips: CPU + memory
- Embedded apps leverage energy efficiency, adjustable mem. capacity, smaller board area
  - $\Rightarrow$  OK market v. desktop (55M 32b RISC '96)

# IRAM 1000

## not a new idea

Stone, '70 "Logic-in memory"

Barron, '78 "Transputer" 100

Dally, '90 "J-machine"

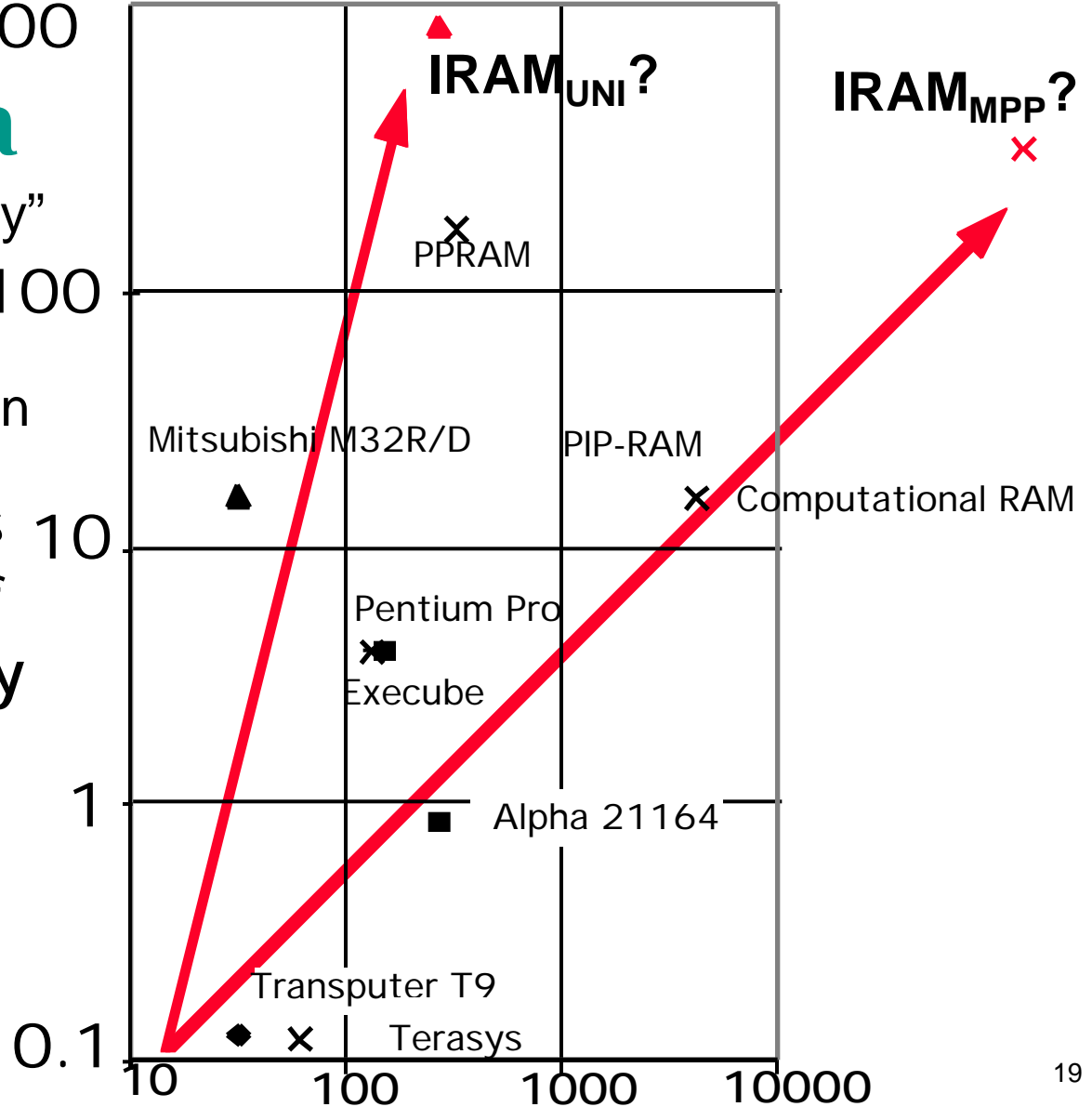
Patterson, '90 panel session

Kogge, '94 "Execube"

Mbits  
of  
Memory

- ✕ SIMD on chip (DRAM)
- Uniprocessor (SRAM)
- ✕ MIMD on chip (DRAM)
- ▲ Uniprocessor (DRAM)
- ◆ MIMD component (SRAM)

Bits of Arithmetic Unit



# Goal for Vector IRAM Generations

- V-IRAM-1 ( $\approx 2000$ )
- 256 Mbit generation (0.20)
- Die size = 256 Mb DRAM die
- 1.5 - 2.0 v logic, 2-10 watts
- 100 - 500 MHz
- 4 64-bit pipes/lanes
- 1-4 GFLOPS(64b)/6-32G (8b)
- 30 - 50 GB/sec Mem. BW
- 24 MB capacity + DRAM bus
- Several fast serial I/O
- V-IRAM-2 ( $\approx 2003$ )
- 1 Gbit generation (0.13)
- Die size = 1 Gb DRAM die
- 1.0 - 1.5 v logic, 2-10 watts
- 200 - 1000 MHz
- 8 64-bit pipes/lanes
- 2-16 GFLOPS/24-128G
- 100 - 200 GB/sec Mem. BW
- 96 MB cap. + DRAM bus
- Many fast serial I/O

# Simple v. Complex Case Study

MIPS MPUs	R5000	R10000	10k/5k
■ Clock Rate	200 MHz	195 MHz	1.0x
■ On-Chip Caches	32K/32K	32K/32K	1.0x
■ Instructions/Cycle	1(+ FP)	4	4.0x
■ Pipe stages	5	5-7	1.2x
■ Model	In-order	Out-of-order	---
■ SPECint_base95	5.7	8.8	1.6x
■ Die Size (mm <sup>2</sup> )	84	298	3.5x
– without cache, TLB	32	205	6.3x
■ Development (man yr.)	60	300	5.0x

# Processor-Memory Performance Gap “Tax”

Processor	% Area ( <i>≈cost</i> )	%Transistors ( <i>≈power</i> )
■ Alpha 21164	37%	77%
■ StrongArm SA110	61%	94%
■ Pentium Pro	64%	88%
– 2 dies per package: Proc/I\$/D\$ + L2\$		
■ Caches have no inherent value, only try to close performance gap		

# How to get Low Power, High Clock rate IRAM?

- Digital Strong ARM 110 (1996): 2.1M Xtors
  - 160 MHz @ 1.5 v = 184 “MIPS” < 0.5 W
  - 215 MHz @ 2.0 v = 245 “MIPS” < 1.0 W
- Start with Alpha 21064 @ 3.5v, 26 W
  - Vdd reduction  $\Rightarrow$  5.3X  $\Rightarrow$  4.9 W
  - Reduce functions  $\Rightarrow$  3.0X  $\Rightarrow$  1.6 W
  - Scale process  $\Rightarrow$  2.0X  $\Rightarrow$  0.8 W
  - Clock load  $\Rightarrow$  1.3X  $\Rightarrow$  0.6 W
  - Clock rate  $\Rightarrow$  1.2X  $\Rightarrow$  0.5 W
- 6/97: 233 MHz, 268 MIPS, 0.36W typ., \$49